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Fractional coloring conjecture for shadow tomography

joint work with:

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Shadow tomography

Goal: given ϱ and a set of observables $P \in S$, estimate $\text{tr}(\varrho P)$.

More precisely, when talking about *shadow tomography*:

Definition (Shadow tomography)

The shadow tomography task for a set S of observables is as follows. We are given copies of an unknown state ϱ , and our goal is to output estimates y_P such that (with high probability) we have $|y_P - \text{tr}(P\varrho)| \leq \epsilon$ for all $P \in S$.

King, et al. Triply efficient shadow tomography. PRX Quantum, 6(1):010336, 2025

→ focus here on Pauli observables

Naive strategy: coloring

Each commuting set can be measured simultaneously.

Formalize: Let G be the anti-commutativity graph of the Pauli set, where two vertices are connected, if and only if the corresponding operators anti-commute

$$\begin{aligned} i \sim j & \quad \text{if } P_i P_j = -P_j P_i \\ i \not\sim j & \quad \text{if } P_i P_j = P_j P_i \end{aligned}$$

- ▶ Color graph with k colors, so that any two adjacent vertices have a different color.
 - ▶ k sets of mutually commuting operators.
 - ▶ The smallest number of colors needed for such coloring is the chromatic number χ . (NP-hard)
- The number of samples for naive tomography scales with χ .

Better strategy: fractional coloring

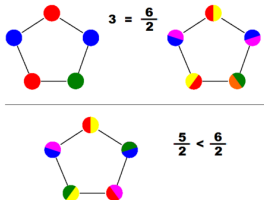
- ▶ probabilistically sample independent sets (sets of mutually commuting elements) with some overlaps.
- ▶ The natural relaxation is the *fractional chromatic number*

The *fractional chromatic number* is a linear-programming relaxation of the chromatic number,

$$\chi_f(G) = \min \left\{ \sum_{I \in \mathcal{I}(G)} w_I : w_I \geq 0, \quad \sum_{I: v \in I} w_I \geq 1 \text{ for every } v \in V(G) \right\}.$$

where $\mathcal{I}(G)$ is the set of independent sets of G .

Each independent set receives a nonnegative weight and every vertex must receive total weight at least one.



Fractional coloring strategy

- ▶ With the fractional coloring strategy, the number of samples scales with the fractional coloring number χ_f .
- ▶ Recall: We actually only care up to precision ϵ :

Definition (Shadow tomography)

The shadow tomography task for a set S of observables is as follows. We are given copies of an unknown state ϱ , and our goal is to output estimates y_P such that (with high probability) we have $|y_P - \text{tr}(P\varrho)| \leq \epsilon$ for all $P \in S$.

King et al showed that the number of samples scales as:

$$\text{samples: } \Omega(\chi_f \log(|S|)/\epsilon^2),$$

$$\text{runtime: } \Omega((T + n^3)\chi_f(\log |S|)/\epsilon^2).$$

Fractional coloring conjecture

In what follows:

$$B_\epsilon(\varrho_m) = \{P \in \mathcal{P}_n : |\text{tr}(\varrho_m P)| \geq \epsilon\},$$

(The set of observables we care about, with sufficient large expectation values.)

Conjecture (Conjecture 13, King et al.)

Let ϱ be an n -qubit state, $\epsilon \in (0, 1)$, and let $B_\epsilon \subseteq \mathcal{P}_n$ be the set of all Paulis P such that $|\text{tr}(\varrho P)| \geq \epsilon$. There is a fractional coloring of the commutation graph $G(B_\epsilon)$ of size $O(1/\epsilon^2)$.

King, et al. Triply efficient shadow tomography. PRX Quantum, 6(1):010336, 2025

- ▶ If true, there would exist a “triply efficient” shadow tomography protocol for the set of all Paulis.
- ▶ We show that this conjecture is false.

Fractional coloring conjecture

Rephrased as an inequality, the Conjecture asserts:
there exists a constant C , such that for every state ϱ and error ϵ

$$\chi_f(B_\epsilon(\varrho)) \cdot \epsilon^2 \leq C,$$

where χ_f is the fractional chromatic number of the anticommutation graph of $B_\epsilon(\varrho)$.

- Thus we need to show that no such finite constant C exists.

Conjecture: There exists a constant C , such that for all states and observables

$$\chi_f(B_\epsilon(\varrho)) \cdot \epsilon^2 \leq C,$$

- ▶ Construct a family of states and observables for which C must necessarily be unbounded.
- ▶ We use the idea of amplification through the lexicographic product. This way the lhs $\rightarrow \infty$.

Lexicographic product

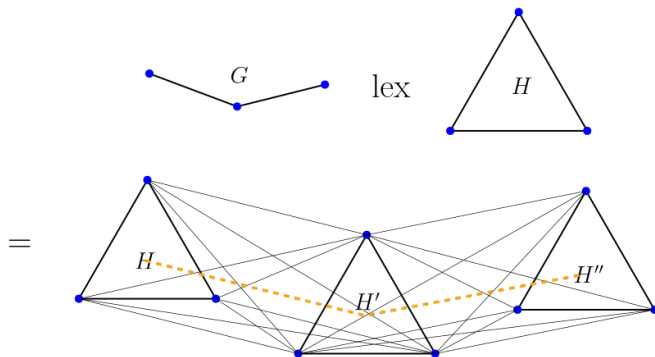
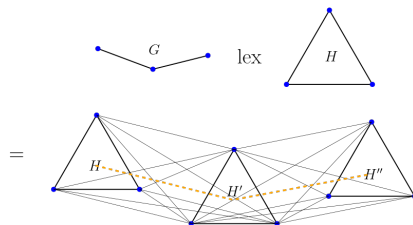


Figure: The lexicographic product of the line graph with the triangle, $L_3 \text{ lex } C_3$. At every vertex of G (a branch), place a copy of H (a leaf). Then connect the vertices among different leaves if and only if their branches are connected.

Lex product of observables

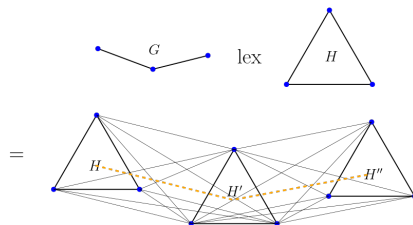


Let $G = L_3$ and $H = C_3$. The graph G is realized by $T = \{XI, ZZ, IX\}$, and H is realized by $R = \{X, Y, Z\}$. A set realizing $(G \text{ lex } H)$ is

leaf H :	$XI XII$,	$XI YII$,	$XI ZII$,
leaf H' :	$ZZ IXI$,	$ZZ IYI$,	$ZZ IZI$,
leaf H'' :	$IX IIX$,	$IX I IY$,	$IX I IZ$.

→ for every copy (leaf) use a separate tensor factor/register

Lex product of observables II



$$T \text{ lex } R =$$

$$\left\{ \overbrace{T_i}^{\text{Reg. 1}} \otimes \overbrace{I \otimes \cdots \otimes I \otimes \underbrace{R_j}_{i'\text{th tensor factor}} \otimes I \otimes \cdots \otimes I}^{\text{Reg. 2}} \mid i = 1, \dots, t; j = 1, \dots, r \right\}.$$

Here T_i is placed on the first register, and R_j placed onto the i -th tensor factor of the second register, where the second register contains $|V(G)|$ internal tensor factors.

Lex product properties

It is known that

$$\begin{aligned}\chi_f(G \text{ lex } H) &= \chi_f(G)\chi_f(H), & (\text{fract. coloring number}) \\ \alpha(G \text{ lex } H) &= \alpha_f(G)\alpha_f(H). & (\text{independence number})\end{aligned}$$

Proof sketch

- Take anti-heptagon. A Pauli realization is:

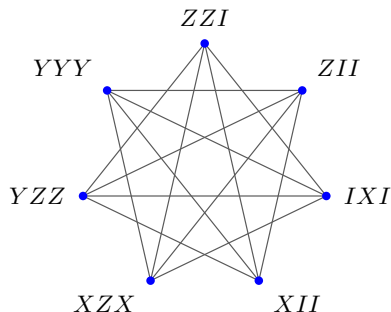


Figure: Anticommutation graph of $\overline{C}_7$ and its Pauli realization.

It has fractional coloring number $\chi_f(\overline{C}_7) = \frac{7}{2}$

(What is its coloring number?)

Proof sketch II

- ▶ Take ground state of Hamiltonian

$$H = \sum_{A \in \mathcal{A}} A$$

Let ϱ be its ground state.

- ▶ One can check that on all observables $A \in \mathcal{A}$ the expectation values are equal,

$$|\mathrm{tr}(A\varrho_{\min})| = \frac{(1 + 2\sqrt{2})}{7} =: a.$$

- ▶ Define lex product observables and states,

$$\mathcal{A}^m = \mathcal{A}^{\mathrm{lex} \, m} = (\mathcal{A} \mathrm{lex} \dots \mathrm{lex} \mathcal{A}) \quad (m \text{ times}).$$

$$\varrho_m = \varrho_{\min}^{\otimes K},$$

- ▶ Check:

$$|\mathrm{tr}(\varrho_m A)| = a^m \quad \text{for all } A \in \mathcal{A}^m.$$

Proof sketch III

► For given m set $\epsilon_m = a^m$.

► Let

$$B_{\epsilon_m}(\varrho_m) = \{P \in \mathcal{P}_n : |\mathrm{tr}(\varrho_m P)| \geq \epsilon_m\},$$

► Realize: $\mathcal{A}^{\mathrm{lex} m}$ is a subset of B_{ϵ_m} . (All elements in $\mathcal{A}^{\mathrm{lex} m}$ have expectation value ϵ^m or larger)

► Anticommutation graph $G(\mathcal{A}^m)$ is an induced subgraph of $G(B_{\epsilon_m})$, and thus $\chi_f(\mathcal{A}^m) \leq \chi_f(B_{\epsilon_m})$.

►

$$\begin{aligned}\chi_f(B_{\epsilon_m}) \cdot \epsilon_m^2 &\geq \chi_f(\mathcal{A}^m) \cdot a^{2m} \\ &= \left(\frac{9 + 4\sqrt{2}}{14}\right)^m \approx (1.046918)^m. \quad (1)\end{aligned}$$

Conclusion

- ▶ Sampling complexity of the shadow tomography protocol by King et al. scales at least with

$$\text{samples: } \Omega(\epsilon^{-2.07598} \log(|S|)/\epsilon^2),$$

$$\text{runtime: } \Omega((T + n^3)\epsilon^{-2.07598}(\log |S|)/\epsilon^2).$$

thus overall at least with $\Omega(1/\epsilon^{4.07598})$.

- ▶ Amplification of properties to construct a counterexample is really useful. See also :

Terence Tao, Amplification, arbitrage, and the tensor power trick <https://terrytao.wordpress.com/2007/09/05/amplification-arbitrage-and-the-tensor-power-trick/>